

# MARKOV CHAINS

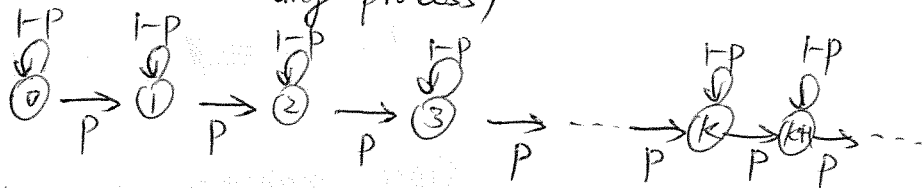
EE464 L25-26-27-28

## Discrete-Time Markov Chains

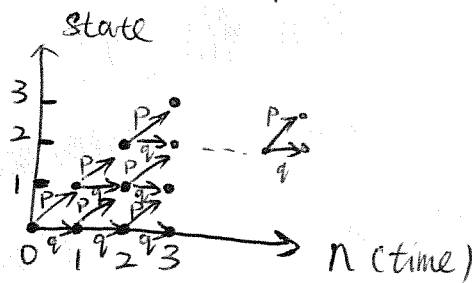
Diagrams =

1. state transition diagram (implicit time)

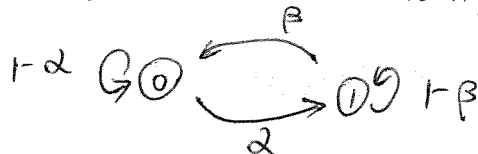
E.g.: (Binomial counting process)



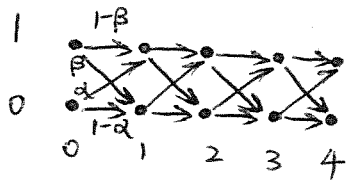
2. Trellis diagram (explicit time)



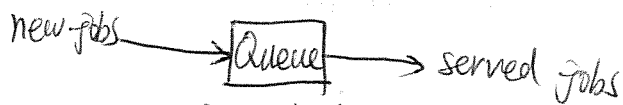
Eg. 2: (2-state Markov chain)



(state transition)



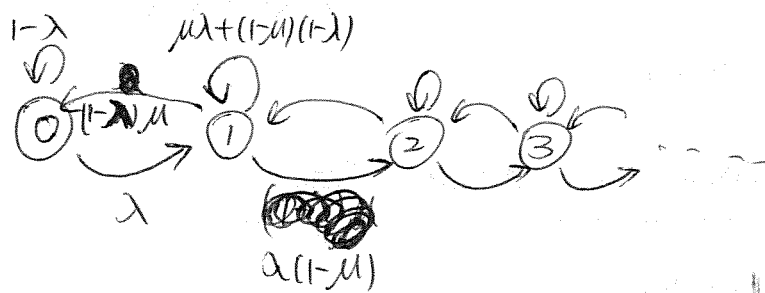
Eg. 3: (Queue):



$X_n$ : # of jobs in queue.

new jobs arrive with prob.  $\lambda$ .

job served w/ prob.  $\mu$ .  $\leftarrow$  iid



Def: Discrete-time Markov chain is homogeneous if

$$P_{ij}^{(n)} \equiv P_R \{ X_{n+1}(u) = j \mid X_n(u) = i \} = P_{ij}$$

Def: State transition probability matrix is

$$P \equiv [P_{ij}]$$

with  $P_{ij} \equiv P_R \{ X_{n+1}(u) = j \mid X_n(u) = i \}$ .

$P$  is a (row) stochastic matrix:

$$\sum_j P_{ij} = 1.$$

Eg. 1: (Binomial counting process)

$$P = \begin{pmatrix} q & p & & & \\ & q & p & & \\ & & q & p & \\ & & & \ddots & \ddots \\ 0 & & & & \ddots \end{pmatrix}_{\infty \times \infty}$$

Eg. 2: (2-state Markov chain)

$$P = \begin{pmatrix} 1-\alpha & \alpha \\ \beta & 1-\beta \end{pmatrix}$$

Eg-3: (Queue)

$$P = \begin{pmatrix} 1-\lambda & \lambda & & & & 0 \\ (1-\lambda)\mu & \mu\lambda + (1-\lambda) & \lambda(1-\mu) & & & \\ & (1-\lambda)\mu & \mu\lambda + (1-\lambda) & \ddots & & \\ & & 0 & \ddots & \ddots & \\ & & & & \ddots & \ddots \end{pmatrix}_{\infty \times \infty}$$

k-step state transition probabilities

Def: k-step state transition prob.

$$P_{ij}^{(k)} \equiv \Pr\{X_{n+k}(u) = j \mid X_n(u) = i\}$$

Thm = (Chapman-Kolmogorov equation)

$$P_{ij}^{(k+l)} = \sum_m P_{im}^{(k)} P_{mj}^{(l)}$$

$\forall k, l \in \mathbb{Z}_+,$

$\forall i, j \in \mathbb{S}$

Proof =  $P_{ij}^{(k+l)} = \Pr\{X_{n+k+l}(u) = j \mid X_n(u) = i\}$

$$= \sum_m \Pr\{X_{n+k+l}(u) = j \mid X_{n+k}(u) = m, X_n(u) = i\}$$

$$= \sum_m \Pr\{X_{n+k}(u) = m \mid X_n(u) = i\} \Pr\{X_{n+l}(u) = j \mid X_n(u) = m\}$$

$$= \sum_m P_{im}^{(k)} P_{mj}^{(l)}$$

State probability vector

Def: state probability vector is a row vector:

$$\Pi(n) \equiv (\Pi_1(n), \Pi_2(n), \dots)$$

where  $\Pi_i(n) \equiv \Pr\{X_n(u) = i\}$

And  $\sum_i \Pi_i(n) = 1.$

Property:

$$\begin{aligned} \Pi_j(n+1) &= \text{Pr}\{X_{n+1}(u) = j\} \\ &= \sum_i \text{Pr}\{X_{n+1}(u) = j, X_n(u) = i\} \\ &= \sum_i \text{Pr}\{X_{n+1}(u) = j \mid X_n(u) = i\} \text{Pr}\{X_n(u) = i\} \\ &= \sum_i P_{ij} \Pi_i(n) \end{aligned}$$

$$\Rightarrow \Pi(n+1) = \Pi(n) P \quad (n=0, 1, 2, \dots)$$

$$\Rightarrow \Pi(n) = \Pi(0) P^n \quad (n=0, 1, 2, \dots)$$

~~n step state transition prob matrix~~

$$\Rightarrow P^n = P^{(n)}$$

Thm = (Matrix version of Chapman-Kolmogorov equation)  
 $P^{(k+l)} = P^{(k)} P^{(l)}$

Proof: This is trivial since  $P^{(n)} = P^n$ .

Def: The steady state, or stationary distribution is

$$\Pi \equiv \lim_{n \rightarrow \infty} \Pi(n)$$

, when the limit exists, and is the same for any  $\Pi(0)$ .

Note that  $\Pi = \Pi P$ . (left Perron vector?)

① Def: Reducible ~~Discrete Time Markov Chain~~ ~~DTMC~~

not all states can be reached from every state.

Note: If for any  $i, j \in S$ , state  $i$  can go to state  $j$   
reducible DTMC can have steady state.

② Def: A state in a DTMC is recurrent if  $P\{ \text{state } i \text{ is revisited} \mid \text{in state } i \} = 1$   
 If a state is not recurrent, we call it's non-recurrent or transient.

Note: 1. recurrent state  $\iff \sum_{m=1}^{\infty} P_{ii}^{(m)} \rightarrow \infty$

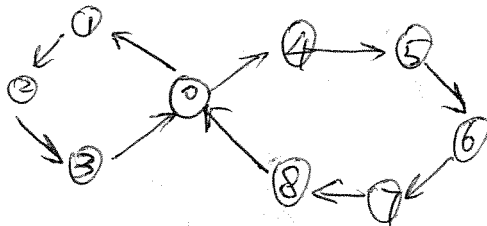
Transient state  $\iff \sum_{m=1}^{\infty} P_{ii}^{(m)} < \infty$

2. ~~An irreducible DTMC has either all recurrent states, or all transient states. (this can only happen for infinite states.)~~ ✓

④ Irreducible, finite state DTMC has all recurrent states.

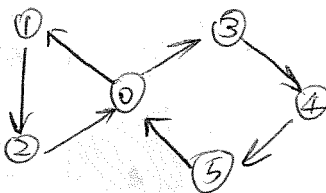
③ Def: A state in a DTMC is periodic w/ period  $d_i > 1$  if

Eg 1°  $P_{ii}^{(m)} = 0, \forall m \neq \text{multiple of } d_i$ .



state 0 is periodic:  
 $d_0 = \text{gcd}(4, 6) = 2$

2°



for state 0:  
 $\text{gcd}(3, 4) = 1$   
 is not periodic.

Note: 1. A DTMC that is irreducible has the same  $d_i$  for every state  $i$ .

①+②+③

Def: An irreducible, aperiodic DTMC is called an ergodic DTMC. <sup>positive recurrent.</sup>

Def: The time to reach state  $j$  <sup>for the first time.</sup> given  $X_0(u) = i$  is a random variable  $T_{ij}(u)$ , with  $1 \leq T_{ij}(u) < \infty$ .

$$m_{ij} \equiv E\{T_{ij}(u)\}$$

~~$m_{ij} \equiv E\{T_{ij}(u)\}$~~

and  $m_{ij}$  is called mean recurrence time -

Thm: 
$$m_{ij} = 1 + \sum_{v \neq j} P_{iv} m_{vj}$$

Proof = Let  $P_{ij}^{*(T_{ij})}$  the prob. from state  $i$  <sup>reaching</sup> state  $j$

for the first time, at time  $T_{ij}$ .

$$\begin{aligned} m_{ij} &= \sum_{T_{ij}=1}^{\infty} T_{ij} P_{ij}^{*(T_{ij})} \\ &= P_{ij} + \sum_{T_{ij}=2}^{\infty} T_{ij} P_{ij}^{*(T_{ij})} \\ &= P_{ij} + \sum_{T_{ij}=2}^{\infty} T_{ij} \sum_{v \neq j} P_{iv} P_{vj}^{*(T_{ij}-1)} \\ &= P_{ij} + \sum_{t=1}^{\infty} (t+1) \sum_{v \neq j} P_{iv} P_{vj}^{*(t)} \\ &= P_{ij} + \sum_{v \neq j} P_{iv} \left( \sum_{t=1}^{\infty} t P_{vj}^{*(t)} + \sum_{t=1}^{\infty} P_{vj}^{*(t)} \right) \end{aligned}$$

$$\Rightarrow \sum_{t=1}^{\infty} P_{ij}^{*(t)} = 1 \quad (\text{irreducibility})$$

$$\sum_{t=1}^{\infty} t P_{ij}^{*(t)} = m_{ij}$$

$$\begin{aligned} \therefore m_{ij} &= P_{ij} + \sum_{v \neq j} P_{iv} (m_{vj} + 1) \\ &= P_{ij} + \sum_{v \neq j} P_{iv} + \sum_{v \neq j} P_{iv} m_{vj} \\ &= 1 + \sum_{v \neq j} P_{iv} m_{vj} \quad \square \end{aligned}$$

Fact: For an irreducible DTMC either

①  $m_{ij} < \infty, \forall i, j \in \mathcal{S} \Rightarrow$  all states are positive recurrent

②  $m_{ii} = \infty, \forall i \in \mathcal{S} \Rightarrow$  all states are null recurrent.

Thm: There are 2 cases for asymptotic behavior of an irreducible, aperiodic DTMC:

(A) the DTMC is <sup>(null recurrent)</sup> not ergodic, then

$P_{ij}^{(n)} \rightarrow 0, \forall i, j$ , there is no stationary  $\Pi$ .

(B) the DTMC is <sup>(positive recurrent)</sup> ergodic, then

•  $\Pi^{(n)} \rightarrow \Pi$ , there is a unique solution to  $\Pi = \Pi P$ .

•  $P_{ij}^{(n)} \rightarrow \Pi_j, \forall j \in \mathcal{S}$ , regardless of  $i$ .

•  $\Pi_j = \frac{1}{m_j}, \forall j \in \mathcal{S}$

• (strong law of large numbers) let  $N_j(u, n)$  is the r.v. <sup>that is</sup> # of times state  $j$  visited in  $\{0, 1, \dots, n\}$ .  $\lim_{n \rightarrow \infty} \frac{N_j(u, n)}{n} = \Pi_j$  Page 4

Methods to find  $\Pi$  :

1. Cut-sets:

Def: Given a DTMC, a cut of  $S$  is a partitioning into  $A, B$ .

Cut-set equations: For an ergodic DTMC, the stationary distribution  $\Pi$  satisfies

$$\sum_{i \in A} \sum_{j \in B} \pi_i P_{ij} = \sum_{k \in B} \sum_{m \in A} \pi_k P_{km}, \quad \forall \text{ cuts } A, B.$$

2. Detailed balance equations:

Thm: For ergodic DTMC, if exists a probability vector  $\alpha$  that satisfies

$$\alpha_i P_{ij} = \alpha_j P_{ji}, \quad \forall i, j \in S$$

then  $\Pi = \alpha$ .

proof =

$$\begin{aligned} \alpha_j &= \alpha_j \sum_i P_{ji} \\ &= \sum_i \alpha_j P_{ji} \\ &= \sum_i \alpha_i P_{ij} \end{aligned}$$

$\therefore \alpha = \alpha P$   
Since the DTMC is ergodic.  
 $\therefore \Pi = \alpha$ .



Eg<sup>o</sup> B/B/1 Queue stationary distribution: ( $\lambda < \mu$ )  
 Bernoulli arrival model, Bernoulli service model, 1 queue

$$\pi_0 = \frac{1}{1 + \frac{\lambda}{\mu \bar{x} - \lambda \bar{u}}}$$

$$\pi_i = \pi_0 \left( \frac{\lambda}{\mu \bar{x}} \right) \cdot \left( \frac{\lambda \bar{u}}{\mu \bar{x}} \right)^{i-1} \quad (i=1, 2, \dots)$$

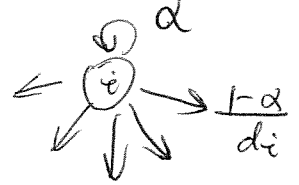
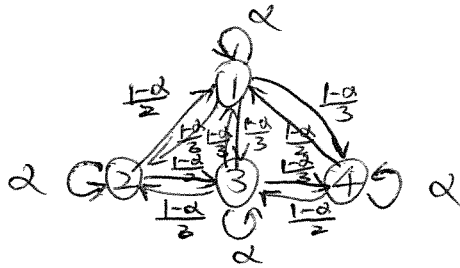
where  $\bar{x} = 1 - \lambda$ ,  $\bar{u} = 1 - \mu$ .

Eg 2<sup>o</sup>: (Random walk over a finite, connected, undirected graph)

~~$X_n(u) \in \{1, 2, \dots, N\}$~~

$X_{n+1}(u) = X_n(u)$  occurs w/ prob.  $\alpha$ .

$X_{n+1}(u)$  is a randomly selected neighbor with equal prob.



~~$d_i = \#$  of neigh~~

$d_i$ : degree of node  $i$  in the original graph.

$$\pi = \frac{1}{\sum_{j=1}^n d_j} (d_1, d_2, \dots, d_n)$$

~~Let  $\alpha_i = C d_i$ ,~~

then  $\alpha_i p_{ij} = C d_i \frac{1-\alpha}{d_i} = C(1-\alpha)$

$\alpha_j p_{ji} = C d_j \frac{1-\alpha}{d_j} = C(1-\alpha)$

$\forall i, j$ .

$\therefore \alpha_i p_{ij} = \alpha_j p_{ji}, \forall i, j$ .

